#### Fine-grained Deterministic Parallelization of Static Analyses

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Joint work with

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# Background

- Associate professor at KTH (2014–2018 assistant professor)
- 2005–2014 Scala language team
  - 2012–2014 Typesafe, Inc. (now Lightbend, Inc.)
- Co-author Scala language specification
- Focus on asynchronous, concurrent and distributed programming
  - Creator of Scala actors, co-author of Scala's futures and async/await
  - *Topics:* programming models, compilers, type systems, semantics

Lightbend

## **The Problem**

- Increasing importance of static analysis
  - Bug finding, security analysis, taint tracking, etc.
- Precise and powerful analyses have *long running times* 
  - Infeasible to integrate into nightly builds, CI, IDE, ...
  - Parallelization difficult: advanced static analyses not data-parallel
- Scaling static analyses to ever-growing software systems requires maximizing utilization of multi-core CPUs

# The Approach

Ongoing work on checking correctness

- Novel concurrent programming model
  - Generalization of futures/promises
  - Guarantees deterministic outcomes (if used correctly)
- Implemented in Scala
  - Statically-typed, integrates functional and object-oriented programming
  - Supported backends: JVM, JavaScript (+ experimental native backend)
- Integrated with OPAL, a state-of-the-art JVM bytecode analysis framework

# Example

- Two key concepts: *cells* and *handlers*
- Cell completers permit *writing*, cells only *reading* (concurrently)

```
val completer1 = CellCompleter[...]
val completer2 = CellCompleter[...]
val cell1 = completer1.cell
val cell2 = completer2.cell
cell2.when(cell1) { update =>
  if (update.value == Impure) FinalOutcome(Impure)
  else NoOutcome
completer1.putFinal(Impure)
```

## Example

- Two key concepts: *cells* and *handlers*
- Cell completers permit *writing*, cells only *reading* (concurrently)

```
val completer1 = CellCompleter[...]
                                                           Α
                                                                          D
val completer2 = CellCompleter[...]
                                                                         cell 2
                                                          cell2
                                                                                cell 1
                                                                 cell1
val cell1 = completer1.cell
                                                                                 imp.
val cell2 = completer2.cell
                                                                          Е
                                                           \mathbf{B}
                                                                         cell 2
                                                                                cell 1
cell2.when(cell1) { update =>
                                                          cell 2
                                                                  cell 1
                                                                         imp
                                                                                 imp.
  if (update.value == Impure) FinalOutcome
  else NoOutcome
                                                                          F
                                                          cell 2
                                                                         cell 2
                                                                  cell 1
                                                                                cell 1
completer1.putFinal(Impure)
                                                                                 imp.
                                                                         imp.
                                                                  imp.
```

#### **Scheduling Strategies**

• **Priorities for message propagations** depending on number of dependencies of source/target nodes and dependees/dependers



#### **Experimental Evaluation**



- Implementation of IFDS<sup>1</sup> analysis framework
- Use IFDS framework to implement taint analysis
  - search for methods in JDK with return type Object or Class with String parameter that is later used in an invocation of Class.forName

<sup>1</sup> Interprocedural Finite Distributive Subset

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Analysis executed on Intel(R) Core(TM) i9-7900X CPU @ 3.30GHz (10 cores) with 128 GB RAM running Ubuntu 18.04.1 and OpenJDK 1.8\_212

### **Scheduling Strategies**

**Table 2.** Performance of different scheduling strategies. Percentages show speedup of each strategy compared to (a) default strategy and (b) slowest strategy.

Strategy	Run	Speedup	Speedup
	time [s]	(a)	(b)
DefaultScheduling	57.15	0.00%	26.54%
SourcesWithManyTargetsLast	53.37	6.62%	31.40%
TargetsWithManySourcesLast	38.94	31.86%	49.94%
TargetsWithManyTargetsLast	77.79	-36.12%	0.00%
SourcesWithManySourcesLast	51.30	10.23%	34.05%

- Using suitable scheduling strategy has big impact on execution time
- Best strategy 49.94% faster than worst strategy, 31.86% faster than default

# Conclusion

- Deterministic concurrent programming model
  - Supporting pluggable, domain-specific scheduling strategies
- Implemented as a library for Scala
- Experimental results for state-of-the-art IFDS-based taint analysis:
  - Speed-up of 4.94x using 10 threads
  - Significant gains using analysis-specific scheduling strategies
- Open-source code available on GitHub: <u>https://github.com/phaller/reactive-async</u>