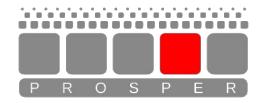
# Formal Verification of Binary Code

Roberto Guanciale



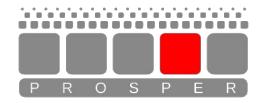
0011111100000011111001001100011011100101	01
11101111110000111010111111011000111101111	10
1101111100001110001101010010111010011111	10
1100011101111011111101111111111010001111	10
11111001101110111100011110111010111010001111	10
11001101111110110001001000000101110111	11
11111000111111000111111111111110011111001111	$\overline{11}$
01111110110011101111011111101011111101111	10
111111111000111111100101001010001111101111	$\overline{01}$
0001111111111101100010100001111001000000	11
1011111001111111101010111111100010111101110000	$\overline{10}$
011100001111011101111100111111100111111	11
1011101101000010011001100011101110000110010000	
1101100110001010111101101000111110100011010	11
11111111100110001111001111010100001101001110001001000	$\overline{10}$
1110100001011100111111100000111111011111	00
1110010011010111111111100111111111111	11
000000001111111100001100111101100100011011011010	01
0100010011111110011111111111111110011111	$\overline{01}$
01000111100100111111100001011110111001101111	11

 $\begin{array}{c} \pi: \{P\} \; \alpha \rightsquigarrow \beta \cup \beta' \{Q\} \\ \\ \overline{\forall \alpha' \in \beta \; .\pi: \{Q\} \; \alpha' \rightsquigarrow \beta' \; \{Q\}} \\ \\ \overline{\pi: \{P\} \; \alpha \rightsquigarrow \beta' \{Q\}} \end{array}$ 



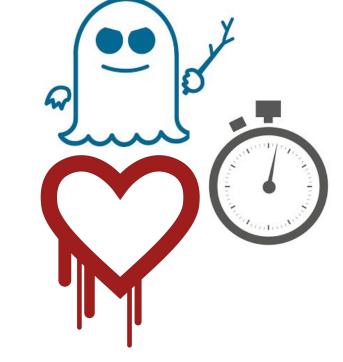


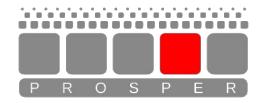




Security. Performance. Proof.



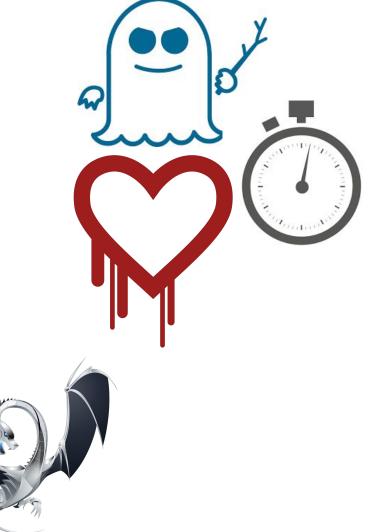


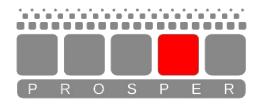


Security. Performance. Proof.











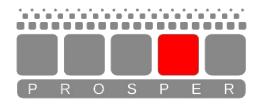


// a0=GETBYTE(s0, 3); ldr r3, [r7, #84] r3, r3, #24 lsrs uxtb r3, r3 r3, [r7, #48] str . . . // v0=\*(Te[0] + a0); ldr r3, [r7, #48] lsls r2, r3, #2 ldr r3, [pc, #928] ; AesEncrypt+0x428 r3, r2, r3 adds ldr r3, [r3, #0] r3, [r7, #32] str . . .

// t0 = v0 ^ v1 ^ v2 ^ v3 ^ rk[0]



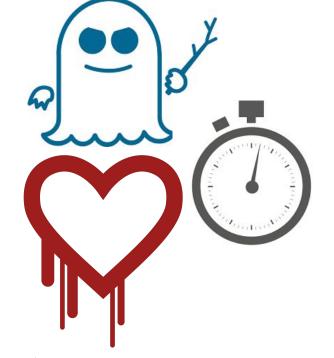








// a0=GETBYTE(s0, 3); ldr r3, [r7, #84] lsrs r3, r3, #24 uxtb r3, r3 r3, [r7, #48] str . . . // v0=\*(Te[0] + a0); ldr r3, [r7, #48] r2, r3, #2 lsls r3 [pc, #928] ; AesEncrypt+0x428 ldr r3, r2, r3 adds r3, [r3, #0] ldr r3, [r7, #32] str . . . // t0 = v0 ^ v1 ^ v2 ^ v3 ^ rk[0]

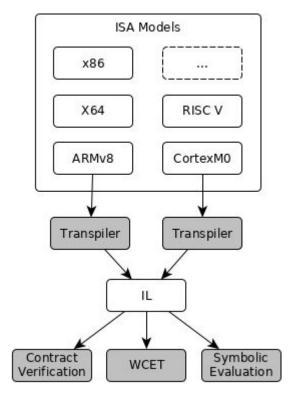




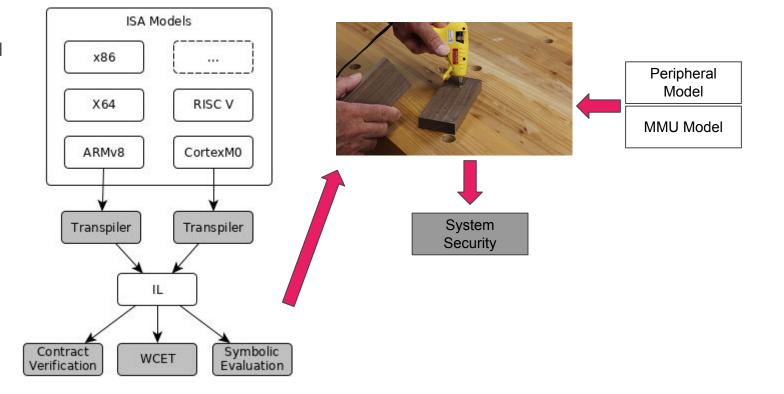
- Valgrind
- BAP

\_ \_\_ \_\_

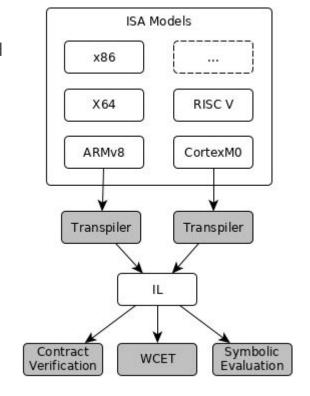
• Angr



- Valgrind
- BAP
- Angr

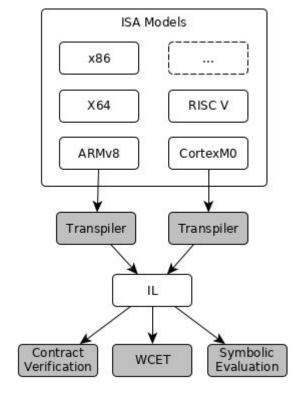


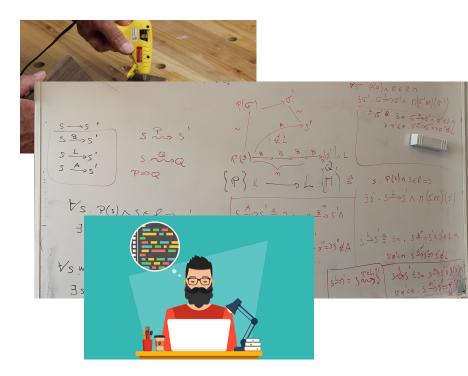
- Valgrind
- BAP
- Angr



S -> S ' S B>S' S Pross' s nong P⇒Q S Los' S As' e \_\_\_\_\_ L(M)  $\frac{\forall s \cdot P(s) \land s \in \ell = s}{\exists m s.' s \cdot \frac{\mathfrak{B}^{m}}{\mathfrak{B}^{m}}} S' \land R(s') \land s' \in L$ SAS' & Jm SB SA S'EA A y'm'<n . S \$55"->5"\$A VS WP(S) ASER-> stos'= 5mgs stos = 3m.s. ∃s'. s → s'AQ(s')

- Valgrind
- BAP
- Angr





# **Certifying (Proof-producing) Analysis of Binaries**

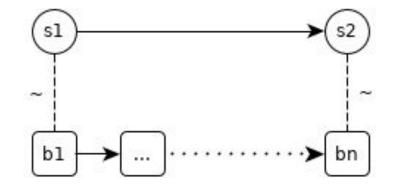
- Implemented using Interactive Theorem Prover (HOL4)
   => Machine checkable proofs
- Formal semantics if ISAs (ARM/Risc-V/etc)
- Formal semantics of BinaryIntermediateRepresentation
  - $\circ$  Similar to LLVM IR
  - $\circ$   $% \left( {{{\rm{Language}}} \right)$  Language designed to automate analysis
    - Program not in memory / Assertions
- Verified theories and proof producing analyses
  - Transpilation
  - Contract based verification
  - 0...

# **Certifying Transpilation**

0: pop R1

4: push R1

```
[0 { R1 := MEM[SP];
    SP := SP-4;
    PC := PC+4;
    JMP 4}]
[4 { MEM := MEM with [SP<-R1];
    SP := SP+4;
    PC := PC+4;
    JMP 8}]
```

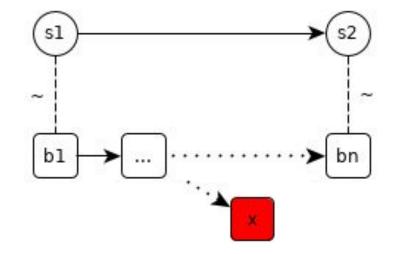


## **Certifying Transpilation**

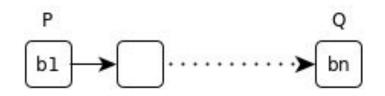
0: pop R1

4: push R1

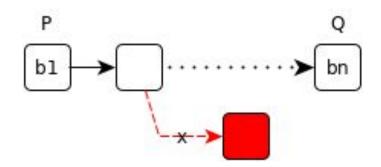
```
[0 { R1 := MEM[SP];
    SP := SP-4;
    PC := PC+4;
    JMP 4}]
[4 { ASSERT(SP not in CODE SECTION);
    MEM := MEM with [SP<-R1];
    SP := SP+4;
    PC := PC+4;
    JMP 8}]
```



- For structured program
  - $\circ \quad \{P\} \text{ statements } \{Q\}$
- For unstructured program?

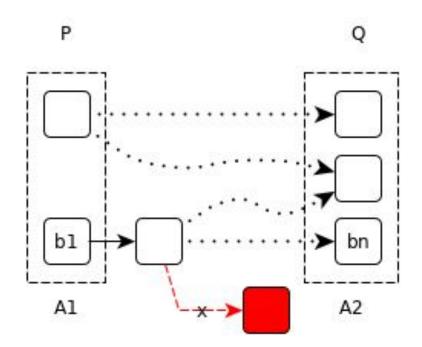


- For structured program
  - $\circ \quad \{P\} \text{ statements } \{Q\}$
- For unstructured program?



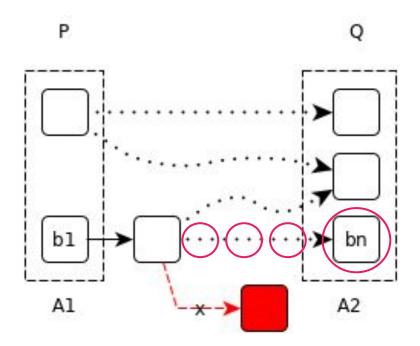
- For structured program

   {P} statements {Q}
   {Q}
- For unstructured program?
  - $\circ$  {P} program: A1 -> A2 {Q}



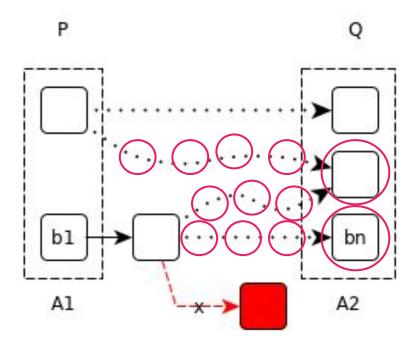
- For structured program

   {P} statements {Q}
   {Q}
- Semi-automatic verification
  - $\circ$   $\,$  Weakest precondition: WP  $\,$
  - $\circ \quad \mathsf{SMT} \text{ solver } \mathsf{P} \Rightarrow \mathsf{WP}$



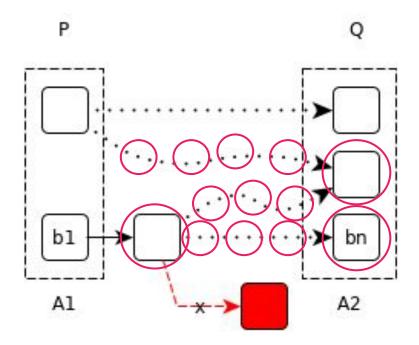
- For structured program

   {P} statements {Q}
   {Q}
- Semi-automatic verification
  - $\circ$   $\,$  Weakest precondition: WP  $\,$
  - $\circ \quad \mathsf{SMT} \text{ solver } \mathsf{P} \Rightarrow \mathsf{WP}$



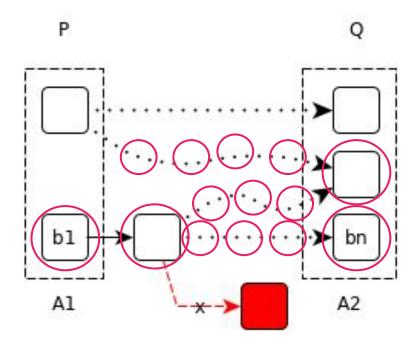
- For structured program

   {P} statements {Q}
   {Q}
- Semi-automatic verification
  - $\circ$   $\,$  Weakest precondition: WP  $\,$
  - $\circ \quad \mathsf{SMT} \text{ solver } \mathsf{P} \Rightarrow \mathsf{WP}$

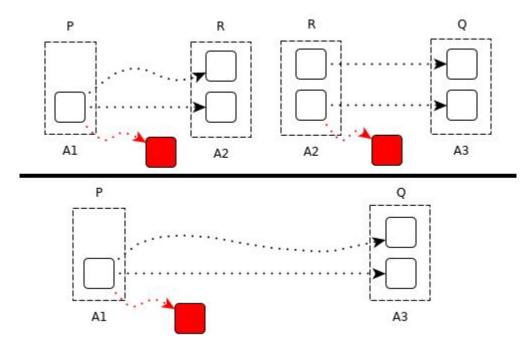


- For structured program

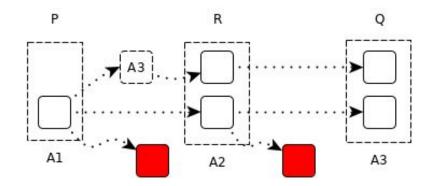
   {P} statements {Q}
   {Q}
- Semi-automatic verification
  - $\circ$   $\,$  Weakest precondition: WP  $\,$
  - $\circ \quad \mathsf{SMT} \text{ solver } \mathsf{P} \Rightarrow \mathsf{WP}$

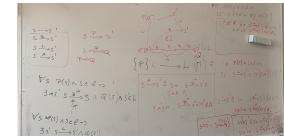


#### **Compositional Logic For Binary Code**

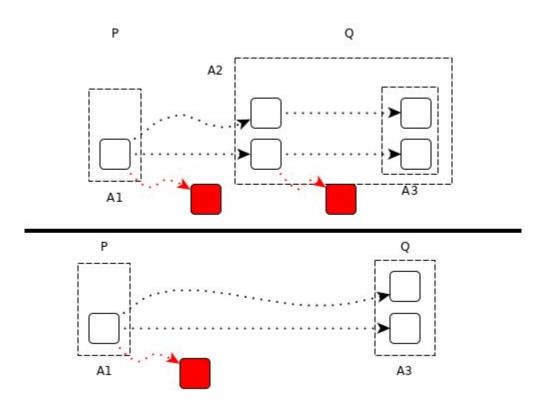


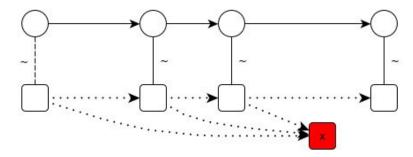
#### **Compositional Logic For Binary Code**



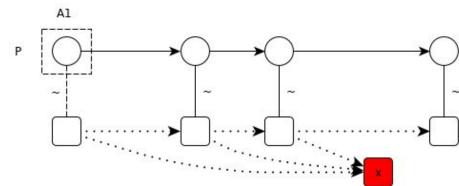


#### **Compositional Logic For Binary Code**

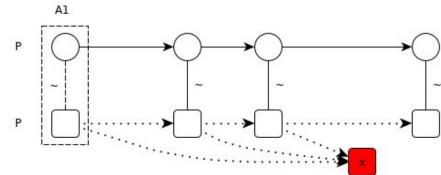




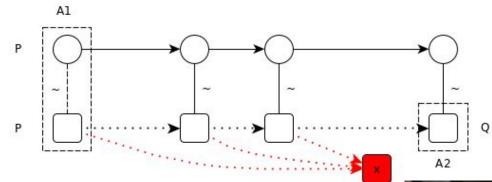




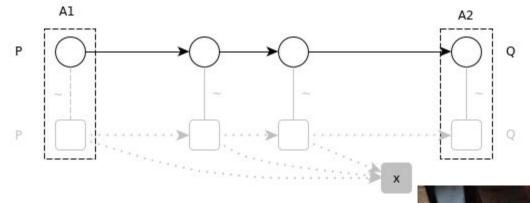














#### Real world usage

#### • Transpilation:

- $\circ$  ~ 5 instructions / s
- $\circ$  numlib / wolf-ssl / lua / SQLite / libc
- $\circ~$  ARMv8 / Cortex M0 / Ongoing Risc-V
- Weakest precondition
  - $\circ$  ~ 1 instruction / s
  - $\circ$  fragments consisting of 10/100 instructions (i.e. AES loop body)



# Thank You

https://github.com/kth-step/HolBA

- Side channel analysis
- Symbolic execution
  - WCET
  - $\circ$  Translation validation
- Kernel verification